



Classical Problem of Synchronization

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Classical Problem of Synchronization

1. Producer Consumer Problem
2. Reader Writer Problem
3. Dining Philosopher Problem
4. Barber Shop Problem

1. Semaphores for Producer Consumer Problem

Buffer size n=5

1

2

3

4

5

Semaphore mutex=1	void producer () { while(true)	void consumer() { while(true)
Semaphore Empty=n (Total number of buffer, suppose n=5)	{ wait (Empty); wait (mutex); append(); signal(mutex); signal(Full);	{ wait(Full); wait(mutex); take(); signal(mutex); signal(Empty);
Semaphore Full=0 (acts as counter to keep track for free buffer).	 } }	 } }

2. Semaphores for Reader Writer Problem

- A database is to be shared among several concurrent processes.
- Some processes may want only to read the database- Readers
- Other may want to update (read + write)-Writers
- Readers and writes are accessing a shared resource by the following rules:
 - Readers can read simultaneously.
 - Only one writer can write at any time.
 - When a writer is writing, no reader can read.
 - If there is any reader reading, all incoming writers must wait. Thus readers have higher priority.

2.Semaphores for Reader Writer Problem (cont.)

	Writer Process	Reader Process
Semaphore mutex=1 Semaphore wrt=1 int readcount=0	do { wait (wrt) write operation signal (wrt) } while (true);	do { wait (mutex) readcount + + if (readcount==1) wait (wrt) signal (mutex) read operation wait (mutex) readcount - - (if readcount == 0) signal(wrt) signal (mutex) } while (true);

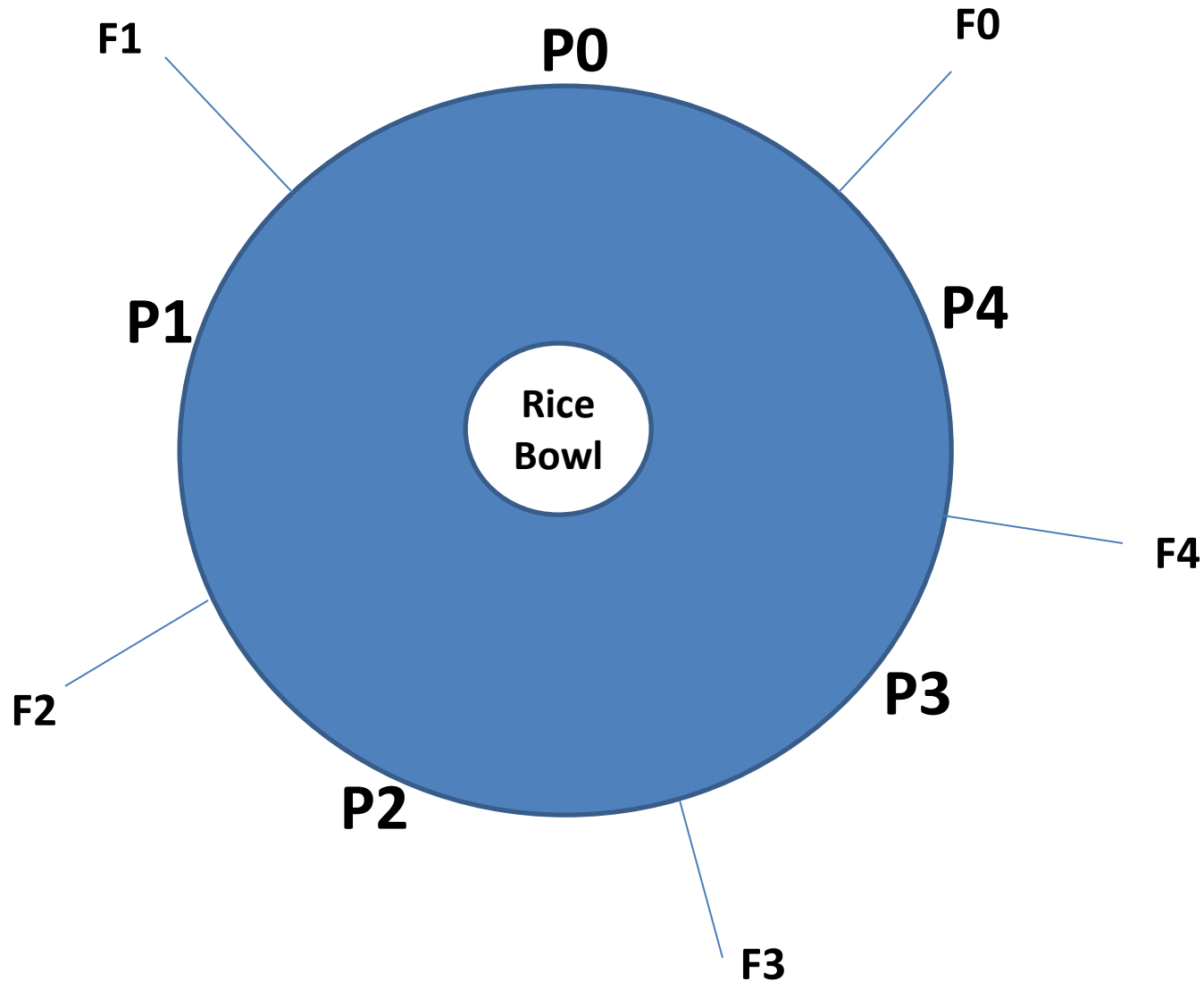
2. Semaphores for Reader Writer Problem (cont.)

- Mutual Exclusion satisfied.
- Progress satisfied.
- But violating bounded waiting: when a write comes in, it waits until no reader is reading.

3. Semaphores for Dining Philosopher Problem

- “Five philosophers sit around a circular table”.
- Each philosopher spend his life alternatively thinking and eating.
- In the centre of the table is a large plate of food.
- A philosopher needs two forks to eat.
- One fork is placed between each pair of philosopher and they agree that each will only use the fork to his immediate left and then right.
- There are five philosopher processes numbered 0 to 4. Fork is also numbered through 0 to 4.

3. Semaphores for Dining Philosopher Problem



3. Semaphores for Dining Philosopher Problem

```
Semaphore s[5];
void philosopher (void)
{
    while(true)
    {
        think();
        wait(takefork(Si)); // Left Fork
        wait(takefork((Si+1)%5)); //Right Fork
        eat();
        signal(putfork(Si));
        signal(putfork((Si+1)%5);
    }
}
```

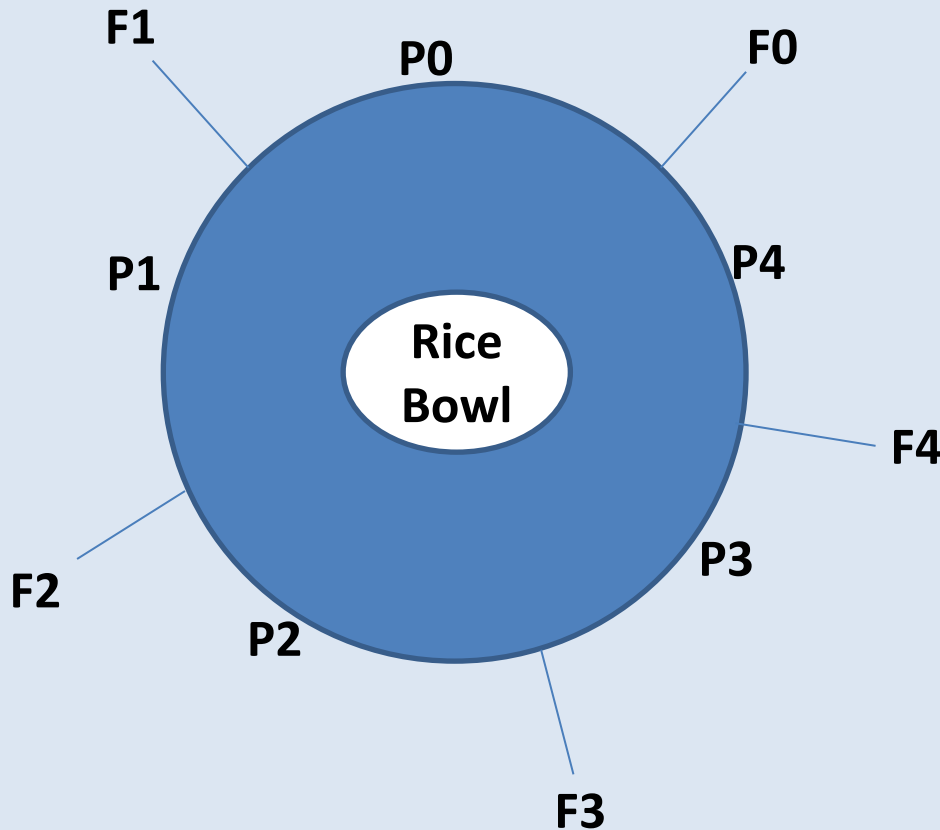
```
s[5]={1,1,1,1,1};
s[0]=1, s[1]=1, s[2]=1,
s[3]=1, s[4]=1
```

P0	S0	S1
P1	S1	S2
P2	S2	S3
P3	S3	S4
P4	S4	S0

Outcome: Deadlock occurred.

3. Semaphores for Dining Philosopher Problem

Solution: P4 Philosopher first take right fork and then take left fork.



$s[0]=1, s[1]=1, s[2]=1, s[3]=1, s[4]=1$

P0	S0	S1
P1	S1	S2
P2	S2	S3
P3	S3	S4
P4	S0	S4

For Nth Process:

wait (takefork($i+1\%5$));

// Right Fork

wait (takefork(i)); //Left Fork

Thank you