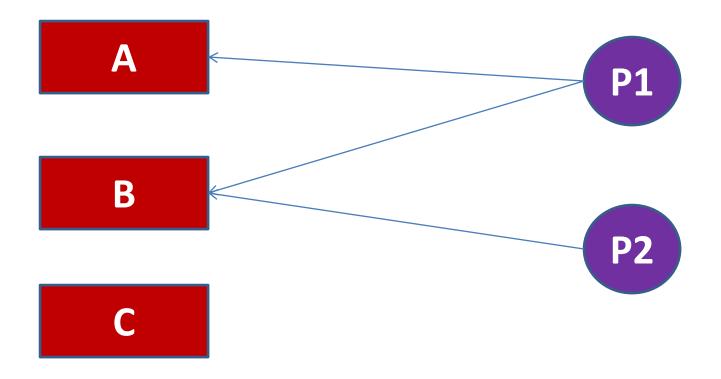
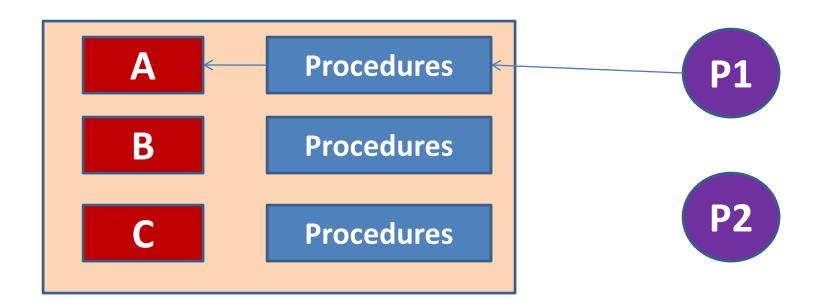




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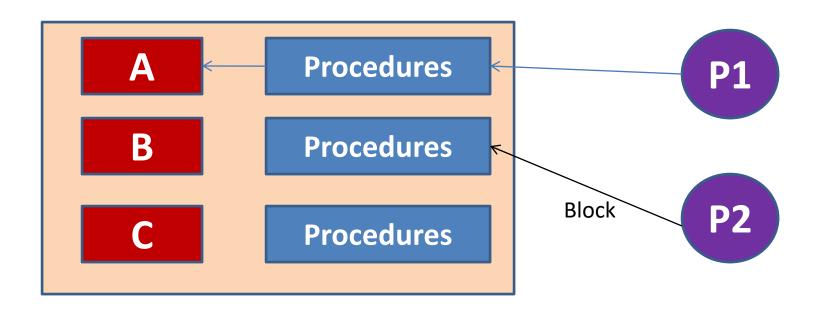


When multiple process access shared resources simultaneously, create problem of race condition.

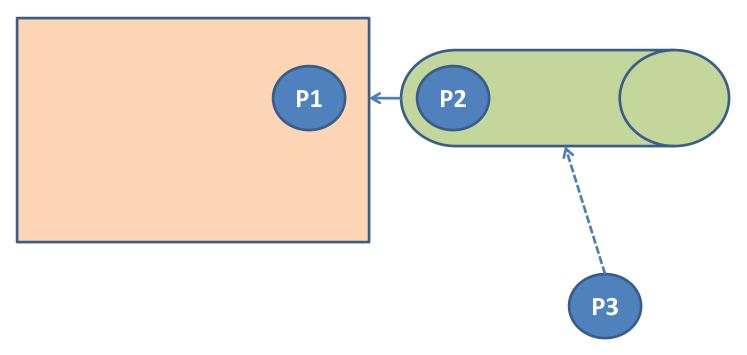


#### A monitor is a module that encapsulates:

- Shared data structure.
- Procedure that operates on shared data.
- Synchronization between concurrent procedure invocation.



Only one process can enter into monitor at a time.



Trying to enter into monitor

```
Monitor account
Double balance;
Withdraw amount()
balance=balance-amount
return balance
```

T1

**T2** 

**T3** 

withdraw (w)

balance=balance-w

Process switch T2

withdraw(a)

Blocked, switch T3

withdraw(b)

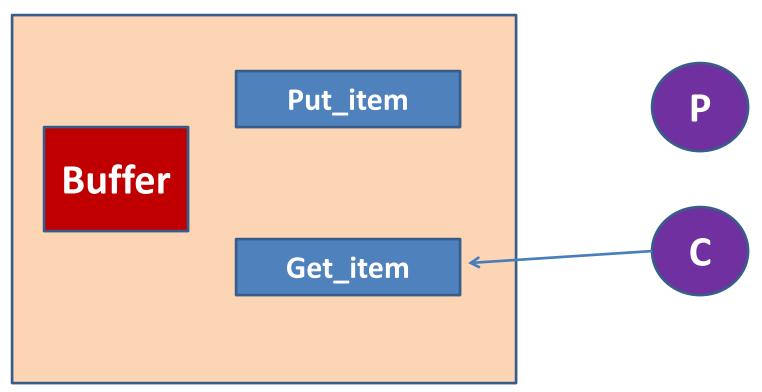
Blocked, switch T1

resume

resume

resume

### **Bounded Buffer Problem**



Initially buffer is empty.

Consumer entered into monitor for consuming data, buffer is empty, it starts waiting in monitor.

Now, producer wants to produce data, but monitor will not allow it to enter because consumer already entered into monitor.

Solution: Conditional Variable

### **Conditional variable**

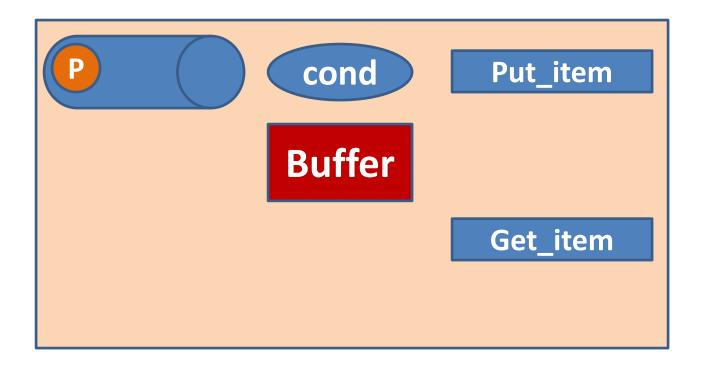
- Conditional variable provides synchronization inside the monitor.
- If a process wants to sleep inside the monitor or it allows a waiting process to continue, in that case conditional variable is used in monitor.
- Three operation can be performed:
  - wait, signal and broadcast.

### **Conditional variable**

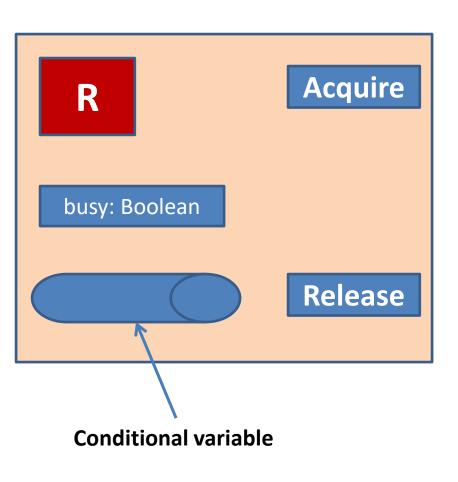
- Wait: if resource is currently not available, current process put to sleep. It releases the lock for monitor.
- Signal: it wakes up one process which are sleeping as a result of wait(). This causes a waiting process to resume immediately. The lock is automatically pass to the waiter, the original process blocked now.
- Broadcast: it signal to all waiting processes.

### **Conditional variable**

Conditionvariable.wait Conditionvariable.signal

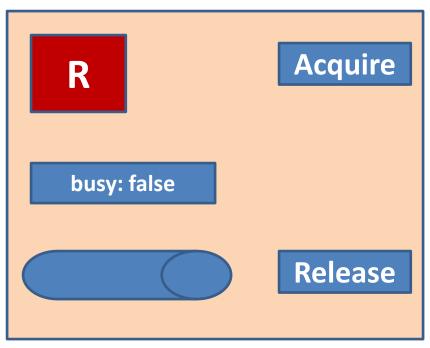


# Single Resource Allocation



```
Monitor single_resource
boolean busy;
condition nonbusy;
  Acquire()
   if busy then nonbusy.wait
   else busy=true
  Release()
   busy=false;
   Nonbusy.signal;
```

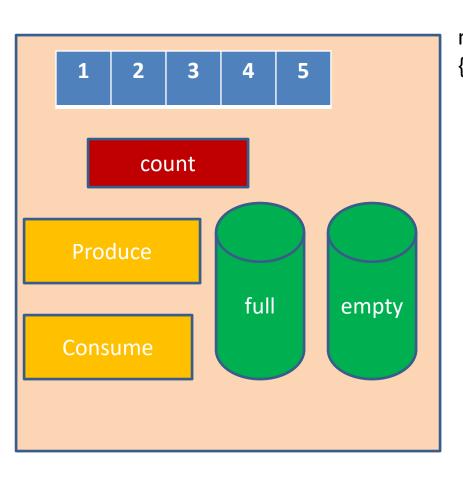
### **Single Resource Allocation**



P1	P2
sr.acquire	sr.acquire
	•••••
	•••••
sr.release	sr.release

**Single Resource** 

#### **Bounded Buffer Problem**



```
monitor bb
 int data[5];
 int count;
 condition full;
 condition empty;
 produce()
  if count==5 then full.wait
  add item to buffer & increase count value
  empty.signal
  consume()
  if count==0 then empty.wait
  access item from buffer and remove count value
  empty.signal
```

# **Thank You**